# Damages and Efficiency with Ex-post Renegotiations

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Lecture 24

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Damages and Reliance

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# ED and Ex-post negotiations I

(Rogerson 1984, RJE)

We assume that

- whenever ex-post negotiations take place, the Buyer gets  $\alpha$  fraction of the *resulting* surplus,
- $\alpha \in [0, 1]$ .  $\alpha$  is exogenously given.

Under Expectation Damages

- in the absence of ex-post negotiations, the breach set  $BS^{ED}(r) = \{C | C > V(r)\}$ , i.e.,  $BS^{ED}(r) = BS^*(r)$ . Therefore,
- there cannot be Pareto improving negotiations.
- As a result, the breach set remains  $BS^{ED}(r) = \{C | C > V(r)\}$ .

Buyer, as before, chooses r that maximizes

$$F(V(r))[V(r) - P] + (1 - F(V(r)))[V(r) - P] - r, i.e.,$$
  
 $V(r) - P - r, i.e.,$ 

the  $r^{ED}$  opted by the Buyer, solves

$$V'(r) - 1 = 0.$$
 (1)

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### RD and Ex-post negotiations I

Under Reliance Damages, in the absence of ex-post negotiations

- the breach set  $BS^{RD}(r) = \{C | C > P + r\}$ , i.e., not Pareto efficient;
- when V(r) > C > P + r holds, the parties can profitably re-negotiate the contract.

Question

- When V(r) > C > P + r, what is the surplus from renegotiations?
- When C > V(r) or C < P + r holds, can there be Pareto improving negotiations?

Therefore, the breach set becomes

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BS^{RD}(r) = \{C|C > V(r)\}.
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The Buyer chooses r that maximizes

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#### RD and Ex-post negotiations II

$$\int_{0}^{P+r} [V(r)-P] dF(C) + \alpha [\int_{P+r}^{V(r)} [V(r)-C] dF(C)] + \int_{P+r}^{\bar{C}} r dF(C) - r, \ i.e.,$$
$$\int_{0}^{P+r} [V(r)-P-r] dF(C) + \alpha [\int_{P+r}^{V(r)} [V(r)-C] dF(C)], \ i.e.,$$

$$(1-\alpha)\int_0^{P+r} [V(r)-P-r]dF(C) + \alpha \int_0^{P+r} [V(r)-P-r]dF(C) + \alpha [\int_{P+r}^{V(r)} [V(r)-C]dF(C)].$$

Rewriting the last two terms, we get

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### RD and Ex-post negotiations III

$$(1-\alpha)\int_{0}^{P+r} [V(r)-P-r]dF(C) + \alpha\int_{0}^{V(r)} \min\left\{\begin{array}{c} V(r)-P-r;\\ V(r)-C \end{array}\right\}dF(C).$$
(2)

When  $\alpha = 0$  the *r* opted by the Buyer, solves

$$V'(r) - 1 = -\frac{f(P+r)[V(r) - P - r]}{F(P+r)}, i.e.,$$
(3)

 $r^{RD}(\alpha = 0)$  opted by the Buyer is the same as in the case of no renegotiations.

Now, from (2) note that the following holds:

- for all  $r \leq r^{ED}$ ,  $V(r) r \leq V(r^{ED}) r^{ED}$
- V(r) and V(r) r both increase with r

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#### • the area of integral increases with r

Therefore, both the terms in (2) and hence their sum will attain a maxima at  $r \ge r^{ED}$ , i.e., for all  $\alpha \in [0, 1]$ ,  $r^{RD}(\alpha) \ge r^{ED}$ .

It can be shown that for all  $\alpha \in [0, 1]$ ,

$$r^{RD}(\alpha = 0) \ge r^{RD}(\alpha) \ge r^{ED}.$$

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## Ex-post negotiations and Relative Efficiency I

#### Proposition

In the presence of ex-post negotiations, Expectation damages are K-H superior to Reliance damages

*Proof*: Consider an arbitrary contract under Reliance damages, say (D(r), P) = (r, P); where *P* is the agreed price.

Suppose, under this contract the outcome is  $(BS^{RD}, r^{RD})$ . We know that for all

$$\alpha \in [0, 1], \ r^{RD}(\alpha) \geq r^{ED}.$$

Also, for any given *r*,

$$Z(r, BS^*(r)) \geq Z(r, BS(r))$$
  
$$Z(r, BS^*(r)) = \int max\{V(r) - r - C, -r\}dF(C).$$

### Ex-post negotiations and Relative Efficiency II

Since  $r^{ED}$  solves max{V(r) - r},

$$V(r^{ED}) - r^{ED} > V(r^{RD}) - r^{RD}$$

Also,  $-r^{ED} > -r^{RD}$ . Therefore

$$Z(r^{\textit{ED}},\textit{BS}^{*\textit{ED}}(r^{\textit{ED}})) \geq Z(r^{\textit{RD}},\textit{BS}^{*}(r^{\textit{RD}})) \geq Z(r^{\textit{RD}},\textit{BS}(r^{\textit{RD}})).$$